

A Platform for Data Intensive Services Enabled by Next Generation Dynamic Optical Networks

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Gesticulation by Franco Travostino



Topics

- Limitations of Packet Switched IP Networks
- Why DWDM-RAM?
- DWDM-RAM Architecture
- An Application Scenario
- Current DWDM-RAM Implementation

Limitations of Packet Switched Networks

What happens when a TeraByte file is sent over the Internet?

- If the network bandwidth is shared with millions of other users, the file transfer task will never be done (World Wide Wait syndrome)
- Inter-ISP SLAs are “as scarce as dragons”
- DoS, route flaps phenomena strike without notice

Fundamental Problems:

- 1) Limited control and isolation of Network Bandwidth*
- 2) Packet switching is not appropriate for data intensive applications => substantial overhead, delays, CapEx, OpEx*

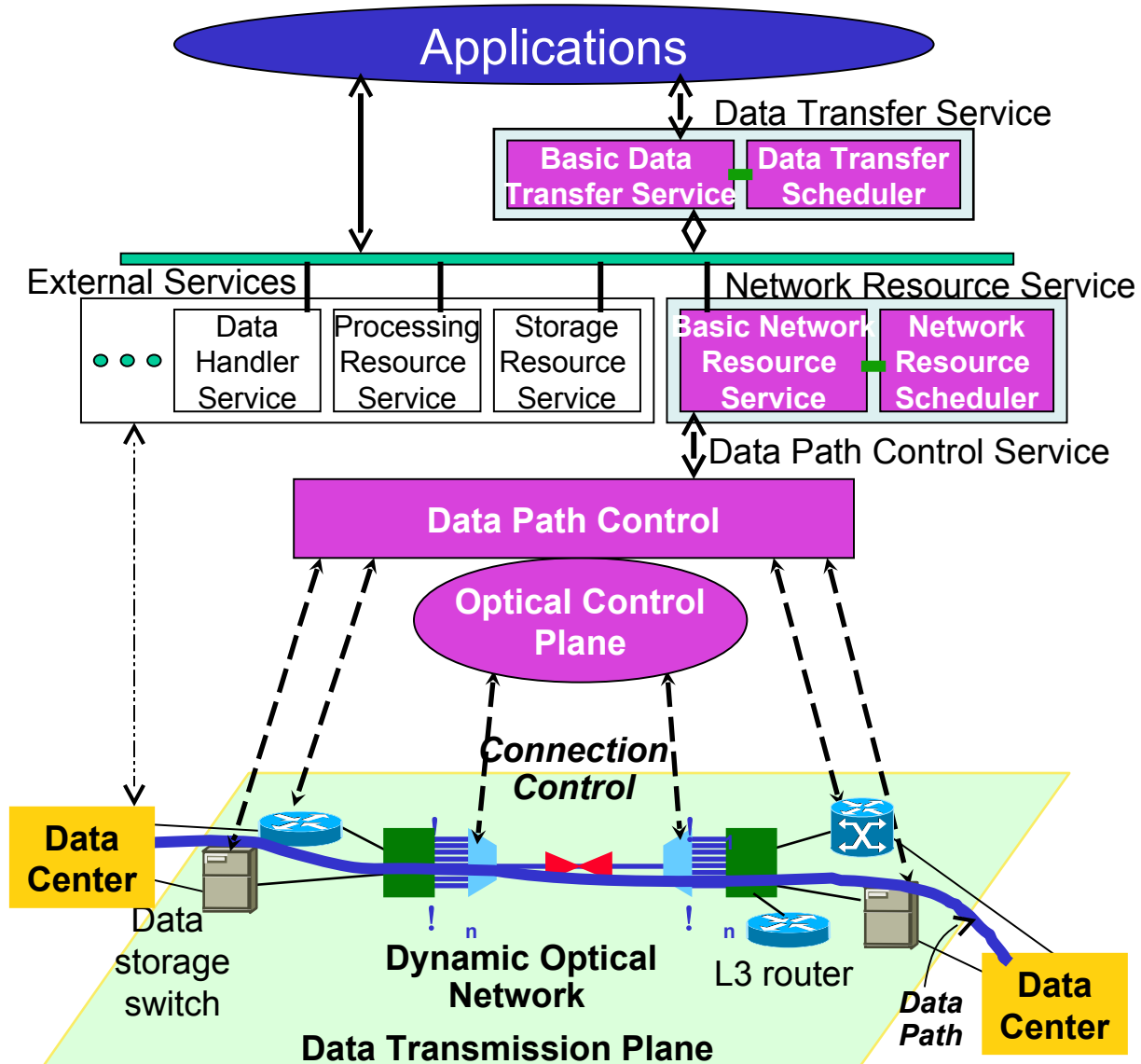
Why DWDM-RAM ?

- The new architecture is proposed for data intensive enabled by next generation dynamic optical networks
 - Encapsulates “optical network resources” into a service framework to support dynamically provisioned and advanced data-intensive transport services
 - Provides a generalized framework for high performance applications over next generation networks, not necessary optical end-to-end
 - Supports both on-demand and scheduled data retrieval
 - Supports a meshed wavelength switched network capable of establishing an end-to-end lightpath in seconds
 - Supports bulk data-transfer facilities using lambda-switched networks
 - Supports out-of-band tools for adaptive placement of data replicas
 - Offers network resources as Grid services for Grid computing

DWDM-RAM Architecture

- The middleware architecture modularizes components into services with well-defined interfaces
- The middleware architecture separates services into 3 principal service layers
 - Data Transfer Service Layer
 - Network Resource Service Layer
 - Data Path Control Service Layer over a Dynamic Optical Network

DWDM-RAM ARCHITECTURE



DWDM-RAM Service Architecture

- The Data Transfer Service (DTS):
 - Presents an interface between an application and a system – receives high-level requests, to transfer named blocks of data
 - Provides Data Transfer Scheduler Service: various models for scheduling, priorities, and event synchronization
- The Network Resource Service (NRS)
 - Provides an abstraction of “communication channels” as a network service
 - Provides an explicit representation of network resources scheduling model
 - Enables capabilities for dynamic on-demand provisioning and advance scheduling
 - Maintains schedules and provisions resources in accordance with the schedule
- Data Path Control Service Layer
 - Presents an interface between the network resource service and the network resources of the underlying network
 - Establishes, controls, and deallocates complete paths across both optical and electronic domains

An Application Scenario: Fixed Bandwidth List Scheduling

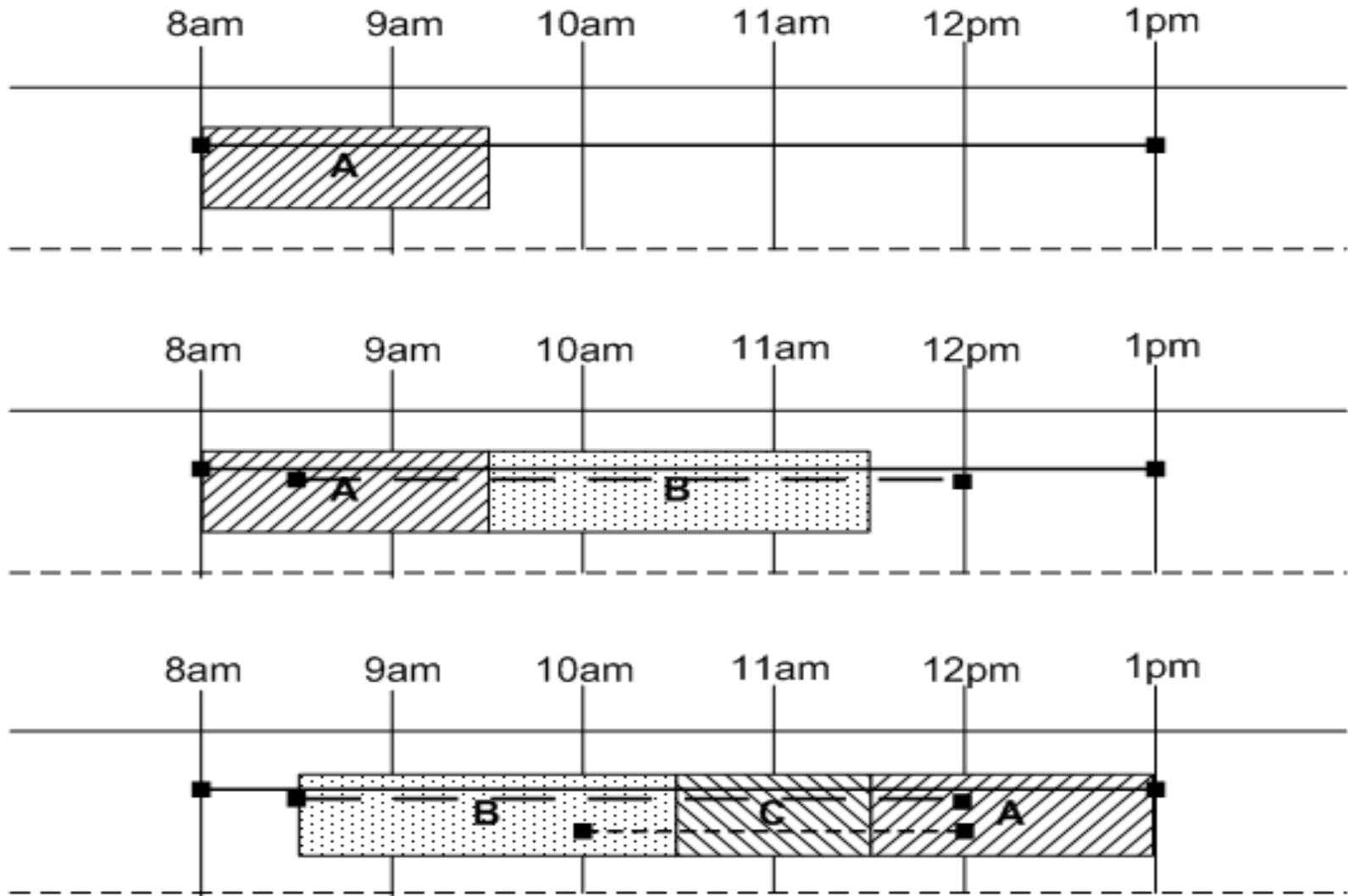
- A scheduling request is sent from the application to the NRS with the following five variables: Source host, Destination host, Duration of connection, Start time of request window, Finish time of request window
- The start and finish times of the request window are the upper and lower limits of when the connection can happen.
- The scheduler must then reserve a continuous hour slot somewhere within that time range. No bandwidth, or capacity, is referred to and the circuit designated to the connection is static.
- The message returned by the NRS is a “ticket” which informs of the success or failure of the request

Fixed Bandwidth List Scheduling

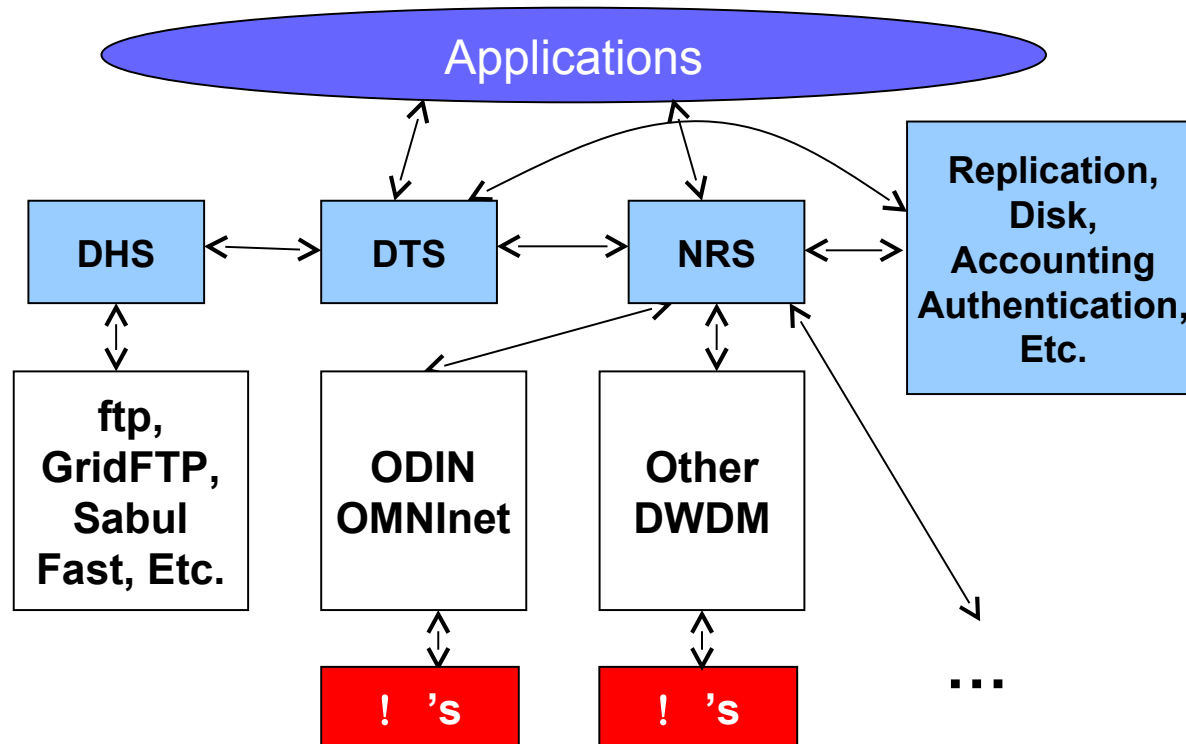
Job	Job Run-time	Window
A	1.5 hours	8am – 1pm
B	2 hours	8:30am – 12pm
C	1 hour	10am – 12pm

This scenario shows three jobs being scheduled sequentially, A, B and C. Job A is initially scheduled to start at the beginning of its under-constrained window. Job B can start after A and still satisfy its limits. Job C is more constrained with its runtime window but is a smaller job. The scheduler adapts to this conflict by intelligently rescheduling each job so all constraints are met.

Fixed Bandwidth List Scheduling



DWDM-RAM Implementation



Dynamic Optical Network

- Gives adequate and uncontested bandwidth to an application's burst
- Employs circuit-switching of large flows of data to avoid overheads in breaking flows into small packets and delays routing
- Is capable of automatic wavelength switching
- Is capable of automatic end-to-end path provisioning
- Provides a set of protocols for managing dynamically provisioned wavelengths

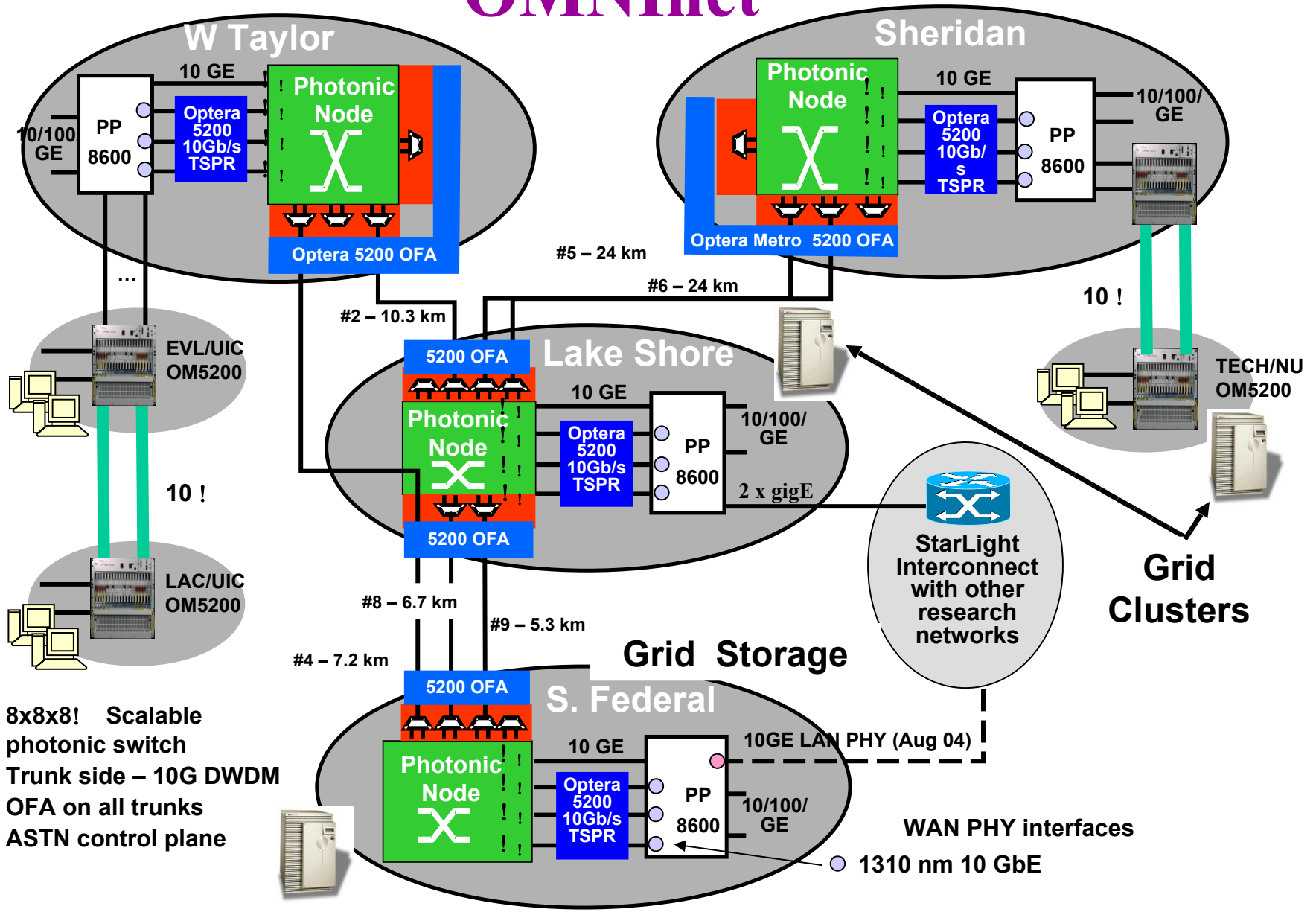
OMNInet Testbed

- Four-node multi-site optical metro testbed network in Chicago -- the first 10GigE service trial when installed in 2001
- Nodes are interconnected as a partial mesh with lightpaths provisioned with DWDM on dedicated fiber.
- Each node includes a MEMs-based WDM photonic switch, Optical Fiber Amplifier (OFA), optical transponders, and high-performance Ethernet switch.
- The switches are configured with four ports capable of supporting 10GigE.
- Application cluster and compute node access is provided by Passport 8600 L2/L3 switches, which are provisioned with 10/100/1000 Ethernet user ports, and a 10GigE LAN port.
- Partners: SBC, Nortel Networks, iCAIR/Northwestern University

Optical Dynamic Intelligent Network Services (ODIN)

- Software suite that controls the OMNInet through lower-level API calls
- Designed for high-performance, long-term flow with flexible and fine grained control
- Stateless server, which includes an API to provide path provisioning and monitoring to the higher layers

OMNIInet

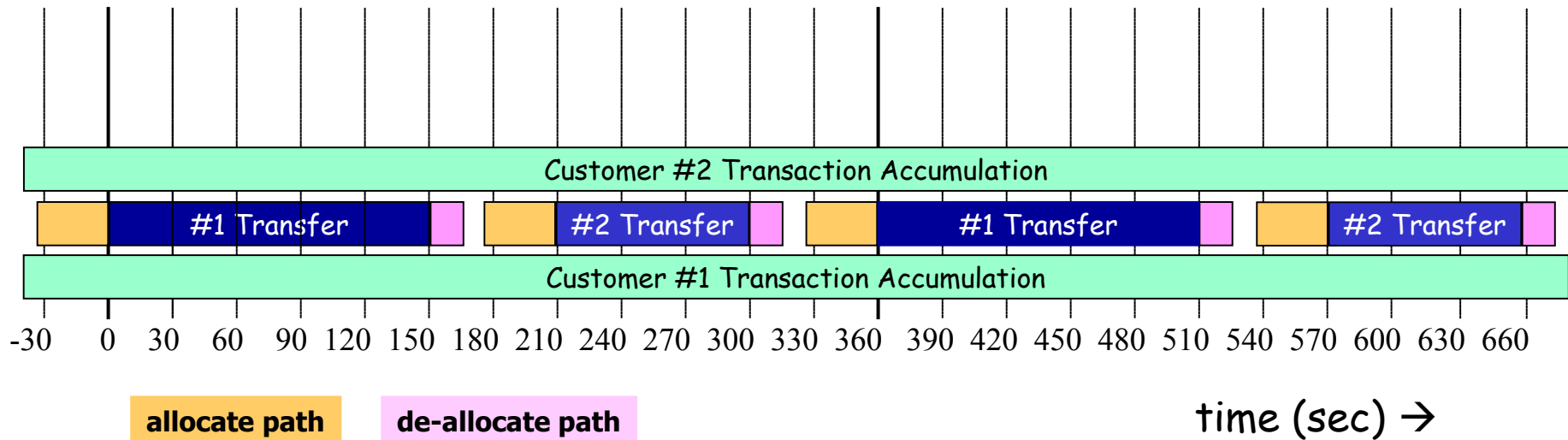


- 8x8x8! Scalable photonic switch
- Trunk side – 10G DWDM
- OFA on all trunks
- ASTN control plane

WAN PHY interfaces
 ○ 1310 nm 10 GbE

Transaction Demonstration Time Line

6 minute cycle time



20GB File Transfer



Conclusion

- The DWDM-RAM architecture yields Data Intensive Services that best exploit Dynamic Optical Networks
- Network resources become actively managed, scheduled services
- This approach maximizes the satisfaction of high-capacity users while yielding good overall utilization of resources
- The service-centric approach is a foundation for new types of services

Some key folks checking us out at our CO2+Grid booth, GlobusWORLD '04, Jan '04



Ian Foster and Carl Kesselman, co-inventors of the Grid (2nd, 5th from the left)

Larry Smarr of OptIPuter fame (6th and last from the left)

Franco, Tal, and Inder (1th, 3rd, and 4th from the left)